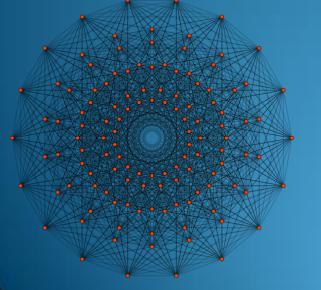


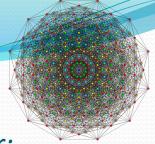
# Partition representable matroids Analysis of paper: Matroid representations by Partitions by Frantisek Matus

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Drexel University Summer 2013 Philadelphia, Pennsylvania







Matroid representations by Partitions by Frantisek Matus

#### General Objective of the Paper:

The aim of this work is to examine partition (p-) representations of matroids along the traditional lines of linear and algebraic representation theories.

#### **Principal Results:**

p-representations of a matroid are closely related to the solutions of a system of generalized quasigroup equations associated to the matroid.

Classify completely the p-representations of a few small matroids and relate them to some classical generalized quasigroup equations.

New examples of non-p-representable matroids are found and an infinite set of excluded minors for the class of p-representable matroids is constructed.



#### Matroid representations by Partitions by Frantisek Matus

#### **Representable Matroids:**

Let  $\mathbb{K}$  be a field. A matroid  $\mathcal{M} = (\mathcal{Q}, \mathcal{I})$  is  $\mathbb{K}$ -representable (or representable for short) if there exists a matrix M over  $\mathbb{K}$  whose columns are indexed by the elements of  $\mathcal{Q}$  such that a subset  $I = \{i_1, ..., i_k\} \subseteq \mathcal{Q}$  is independent if and only if the matrix corresponding columns of M are independent. In this situation, we say that the M is a  $\mathbb{K}$ -representation of the matroid  $\mathcal{M}$ 

#### Linear Matroid:

if  $(E,\mathcal{I})$  is any matroid, then a representation of  $(E,\mathcal{I})$  may be defined as a function f that maps E to a vector space V, with the property that a subset A of E is independent if and only if f(A) is linearly independent.

if V is a vector space over field F then the matroid is called an F-linear matroid.

Thus, the linear matroids are exactly the matroids that are isomorphic to the matroids defined from sets or multisets of vectors. If a matroid is linear, it may be representable over some but not all fields.

Regular Matroids: The unimodular or regular matroids are the matroids that can be represented over all fields; If a matroid is regular, so is its dual matroid, and so is every one of its minors. Every direct sum of regular matroids remains regular. Every graphic matroid (and every co-graphic matroid) is regular. Conversely, every regular matroid may be constructed by combining graphic matroids, co-graphic matroids, and a certain ten-element matroid that is neither graphic nor co-graphic, using an operation for combining matroids that generalizes the clique-sum operation on graphs.

Uniform Matroids: A uniform matroid  $U_n^r$  has n elements, and its independent sets consist of all subsets of up to r of the elements.

Uniform matroids may be represented by sets of vectors in general position in an r-dimensional vector space. The field of representation must be large enough for there to exist n vectors in general position in this vector space, so uniform matroids are F-linear for all but finitely many fields F.



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#### **Partition Matroids:**

Let  $B_i$  be a collection of disjoint sets, and let  $d_i$  be integers with  $0 \le d_i \le |B_i|$ . Define a set I to be "independent" when, for every index i,  $|I \cap B_i| \le d_i$ . Then the sets that are independent sets in this way form the independent sets of a matroid, called a partition matroid. The sets  $B_i$  are called the blocks of the partition matroid. A basis of the matroid is a set whose intersection with every block  $B_i$  has size exactly  $d_i$ , and a circuit of the matroid is a subset of a single block  $B_i$  with size exactly  $d_i + 1$ . The rank of the matroid is  $\sum d_i$ .

Every uniform matroid  $U_n^r$  is a partition matroid, with a single block  $B_1$  of n elements and with  $d_1=r$ . Every partition matroid is the direct sum of a collection of uniform matroids, one for each of its blocks.

#### Transversal Matroids:

In some publications, the notion of a partition matroid is defined more restrictively, with every  $d_i=1$ . The partitions that obey this more restrictive definition are the transversal matroids of the family of disjoint sets given by their blocks.

#### **Graphic Matroids:**

A graphic matroid is the matroid defined from the edges of an undirected graph by defining a set of edges to be independent if and only if it does not contain a cycle. Every graphic matroid is regular, and thus is F-linear for every field F.

#### Gammoids:

Like uniform matroids and partition matroids, the gammoids, matroids representing reachability in directed graphs, are linear over every sufficiently large field. More specifically, a gammoid with n elements may be represented over every field that has at least  $2^n$  elements.



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#### **Algebraic Matroids:**

The algebraic matroids are matroids defined from sets of elements of a field extension using the notion of algebraic independence. Every linear matroid is algebraic, and for fields of characteristic zero (such as the real numbers) linear and algebraic matroids coincide, but for other fields there may exist algebraic matroids that are not linear.

Let *L* be a field. A **subfield** of *L* is a subset *K* of *L* which is closed under the field operations of *L* and under taking inverses in *L*. In other words, *K* is a field with respect to the field operations inherited from *L*. The larger field *L* is then said to be an **extension field** of *K*. In abstract algebra, a subset *S* of a field *L* is **algebraically independent** over a subfield *K* if the elements of *S* do not satisfy any non-trivial polynomial equation with coefficients in *K*.

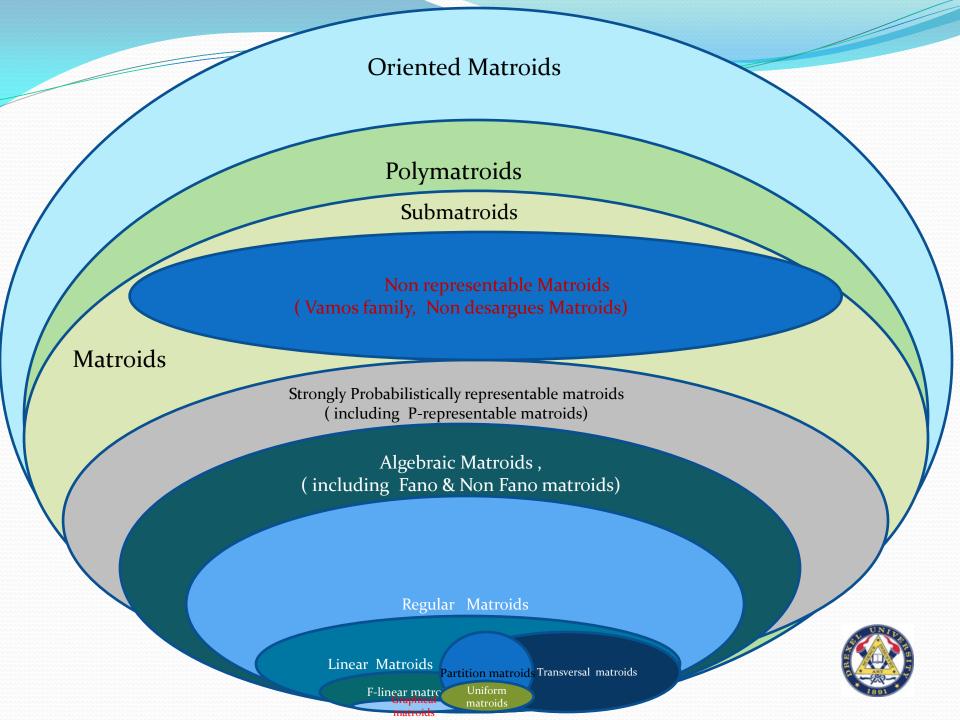
#### Matroids of Entropic Regions:

The **Shannon entropies** of all subvectors of a random vector are considered for the **coordinates of an entropic point** in a Euclidean space. The entropic points in a space define an **entropic region**. The problem to describe the region will be reviewed and connections to the matroid theory discussed.

Two classes of matroids are closely related:

Partition representable and asymptotically entropic. A matroid is **Partition Representable** if its rank function is a multiple of an entropic point. A matroid is **Asymptotically Entropic** if its rank function is in the closure of the entropic region.





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#### Non Representable Matroids:

The first matroid that was known to be not representable was the Vamos Matroid, this structure define an entire family of Non representable matroids. In 2009 the research team of Information Security Lab, College of Computer Science & Technology, Huazhong University of Science and Technology, Wuhan, China, leaded by Chingfang Hsu and Qi Cheng have shown, based on discretization of the Vamos matroids family, that there is a way to generalize a sufficient condition for not representability of matroids.

**Definition** The Vamos matroid is defined on  $Q = \{1, 2, 3, 4, 5, 6, 7, 8\}$  with bases all 4-sets except the five 4-sets which are:  $\{1, 2, 3, 4\}$ ,  $\{1, 2, 5, 6\}$ ,  $\{1, 2, 7, 8\}$ ,  $\{3, 4, 5, 6\}$ ,  $\{3, 4, 7, 8\}$ .

#### Vamos Family

For the Vamos matroid  $\mathcal{M}=(\mathcal{Q},\mathcal{I})$ , there exists a partition  $\Pi_0=\left\{P_1,P_2,P_3,P_4\right\}$  ( $P_1=\left\{1,2\right\},P_2=\left\{3,4\right\},P_3=\left\{5,6,\right\},P_4=\left\{7,8\right\}$ ) of the ground set  $\mathcal{Q}$ , and the partition  $\Pi_0$  defines a mapping  $\Pi_0:\mathcal{P}(\mathcal{Q})\to\mathbb{Z}_+^4$  and, hence, we obtain a discrete polymatroid  $D_V=\Pi_0(\mathcal{I})$  corresponding to the Vamos matroid.



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**Proposition** For a discrete polymatroid D with ground set  $J_m$ , if there exists  $X\subseteq J_m$ , where |X|=4, such that  $D(X)=D_V$ , then D must be a non-representable discrete polymatroid, and hence, the multipartite matroid corresponding to D must be non-representable. All of these discrete polymatroids construct a family of non-representable matroids, that is,  $F_{D_V}=\left\{D\subset \mathbb{Z}_+^m: D(X)=D_V, X\subset J_m\boxplus |X|=4\right\}$ , which we call Vamos Family.

#### A Sufficient Condition for a Discrete Polymatroid to Be Non-representable

**Theorem** Let  $D \subset \mathbb{Z}_+^m$  be a discrete polymatroid with ground set  $J_m$ , if there exists  $X \subseteq J_m$  such that  $D(X) = \{u(X) : u \in D\} \subset \mathbb{Z}_+^{|X|}$  is a non-representable discrete polymatroid, then D must be a non-representable discrete polymatroid and, hence, the multipartite matroid corresponding to D must be non-representable.



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Extension of this method to determine Representability using Secret Sharing schemes in general:

Definition of an Ideal Secrete Scheme from a representable matroid, general procedure:

Let  $\mathbb{K}$  be a finite field and let  $\mathcal{M}=(\mathcal{Q},\mathcal{I})$  be a  $\mathbb{K}$ -representable matroid. Let  $p_0\in\mathcal{Q}$  be special participant called dealer and  $\mathcal{Q}=P\cup\{p_0\}$ . For every  $k\times(n+1)$  matrix M representing  $\mathcal{M}$  over  $\mathbb{K}$ , let E be a vector space of finite dimention  $\dim E=k$  over  $\mathbb{K}$ . For every  $i\in\mathcal{Q}$ , we define a surjective linear mapping:  $\pi_i:E\to\mathbb{K}$ , and the i-th column of M corresponds to the linear form  $\pi_i$ . In that situation, for every random choice of an element  $x\in E$ , we can obtain  $s_i=\pi_i(x)\in\mathbb{K}$  is the share of the participant  $i\in P$  and  $s=\pi_{p_0}(x)\in\mathbb{K}$  is the shared secret value. Hence, by the columns of M, we define an ideal secret sharing scheme with access structure  $\Gamma_{p_0}(\mathcal{M})$ , where  $\min\Gamma_{p_0}(\mathcal{M})=\{A\subseteq P:A\cup\{p_0\}\text{ is a circuit of }\mathcal{M}\}$ . Therefore, the access structures induced by representable matroids are ideal.

Discretization of the Ideal Sharing scheme via Partition of disjoint subsets:

We write  $\mathcal{P}(P)$  for the power set of the set P. An m-partition  $\Pi = \{P_1, ..., P_m\}$  of a set P is a disjoint family of m nonempty subsets of P with  $P = P_1 \cup ... \cup P_m$ . Let  $\Lambda \subseteq \mathcal{P}(P)$  be a family of subsets of P. For a permutation  $\sigma$  on P, we define  $\sigma(\Lambda) = \{\sigma(A) : A \in \Lambda\} \subseteq \mathcal{P}(P)$ . A family of subsets  $\Lambda \subseteq \mathcal{P}(P)$  is said to be  $\Pi$ -partite if  $\sigma(\Lambda) = \Lambda$  for every permutation  $\sigma$  such that  $\sigma(P_i) = P_i$  for every  $P_i \in \Pi$ . We say that  $\Lambda$  is m-partite if it is  $\Pi$ -partite for some m-partition  $\Pi$ . These concepts can be applied to access structures and matroids.



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Construction of a discrete Polymatroid from the Partitioned scheme, in the most general case.

For every integer  $m \geq 1$ , we consider the set  $J_m = \{1, ..., m\}$ . Let  $\mathbb{Z}_+^m$  denote the set of vectors  $u = (u_1, ..., u_m) \in \mathbb{Z}^m$  with  $u_i \geq 0$  for every  $i \in J_m$ . For a partition  $\Pi = \{P_1, ..., P_m\}$  of a set P and for every  $A \subseteq P$  and  $i \in J_m$ , we define  $\Pi_i(A) = |A \cap P_i|$ . Then the partition  $\Pi$  defines a mapping  $\Pi : \mathcal{P}(P) \to \mathbb{Z}_+^m$  by considering  $\Pi(A) = (\Pi_1(A), ..., \Pi_m(A))$ . If  $\Lambda \subseteq \mathcal{P}(P)$  is  $\Pi$ -partite, then  $A \in \Lambda$  if and only if  $\Pi(A) \in \Pi(\Lambda)$ . That is,  $\Lambda$  is completely determined by the partition  $\Pi$  and the set of vectors  $\Pi(\Lambda) \subset \mathbb{Z}_+^m$ .

#### Definition of needed operators on the vector spaces :

If  $u,v\in\mathbb{Z}_+^m$ , we write  $u\le v$  if  $u_i\le v_i$  for every  $i\in J_m$ , and we write u< v if  $u\le v$  and  $u\ne v$ . The vector  $w=u\lor v$  is defined by  $w_i=\max(u_i,v_i)$ . The modulus of a vector  $u\in\mathbb{Z}_+^m$  is  $|u|=u_1+\dots+u_m$ . For every subset  $X\subseteq J_m$ , we write  $u(X)=(u_i)_{i\in X}\in\mathbb{Z}_+^{|X|}$  and  $|u(X)|=\sum_{i\in X}u_i$ 

#### Conceptualization of the Discrete Polymatroid:

A discrete polymatroid on the ground set  $J_m$  is a nonempty finite set of vectors  $D \subset \mathbb{Z}_+^m$  satisfying:

- 1. if  $u \in D$  and  $v \in \mathbb{Z}_+^m$  is such that  $v \le u$ , then  $v \in D$ , and
- 2. for every pair of vectors  $u, v \in D$  with |u| < |v|, there exists  $w \in D$  with  $u < w \le u \lor v$ .



#### Matroid representations by Partitions by Frantisek Matus

A **relation between multipartite matroids and discrete polymatroids** emerge directly from the axioms of independent sets and can be expressed by the following proposition:

**Proposition 2.1.** Let  $\Pi$  be a partition of a set  $\mathcal{Q}$  and let  $\mathcal{I} \subseteq \mathcal{P}(\mathcal{Q})$  be a  $\Pi$ -partite family of subsets. Then  $\mathcal{I}$  is the family of the independent sets of a  $\Pi$ -partite matroid  $\mathcal{M} = (\mathcal{Q}, \mathcal{I})$  if and only if  $\Pi(\mathcal{I}) \subset \mathbb{Z}_+^m$  is a discrete polymatroid.

The bases of the Polymatroid are defined as:

**Proposition 2.2.** A nonempty subset  $\mathcal{B} \subset \mathbb{Z}_+^m$  is the family of bases of a discrete polymatroid if and only if it satisfies:

- 1. all elements in  $\mathcal{B}$  have the same modulus, and
- 2. for every  $u \in \mathcal{B}$  and  $v \in \mathcal{B}$  with  $u_i > v_i$ , there exists  $j \in J_m$  such that  $u_j < v_j$  and

 $u-e_i+e_j\in\mathcal{B}$  , where  $e_i$  denotes the i-th vector of the canonical basis of  $\mathbb{Z}^m$  .

While its rank function is given by:

The rank function of a discrete polymatroid D with ground set  $J_{\it m}$  is the function

 $h: \mathcal{P}(J_m) \to \mathbb{Z}$  defined by  $h(X) = \max\{|u(X)|: u \in D\}$ . The next proposition is a consequence of Properties of its rank function are:

**Proposition 2.3.** A function  $h: \mathcal{P}(J_m) \to \mathbb{Z}$  is the rank function of a discrete polymatroid with ground set  $J_m$  if and only if it satisfies

- 1.  $h(\phi) = 0$ , and
- 2. h is monotone increasing: if  $X \subseteq Y \subseteq J_m$ , then  $h(X) \le h(Y)$ , and
- 3. h is submodular: if  $X, Y \subseteq J_m$ , then  $h(X \cup Y) + h(X \cap Y) \le h(X) + h(Y)$ .



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Characterization of the Polymatroid using its rank function:

$$D = \left\{ u \in \mathbb{Z}_+^m : |u(X)| \le h(X) \text{ for all } X \subseteq J_m \right\}$$

Definition of representability of Matroid based on its rank function:

Let  $\mathbb{K}$  be a field, E a  $\mathbb{K}$ -vector space, and  $V_1,...,V_m$  subspaces of E. It is not difficult to check that the mapping  $h: \mathcal{P}(J_m) \to \mathbb{Z}$  defined by  $h(X) = \dim(\sum_{i \in X} V_i)$  is the rank function of a discrete polymatroid  $D \subset \mathbb{Z}_+^m$ . In this situation, we say that D is  $\mathbb{K}$ -representable and the subspaces  $V_1,...,V_m$  are a  $\mathbb{K}$ -representation of D.

Criteria of Representability of the Secret Shared scheme Matroid based in the associated discrete polymatroid built through partitions:

**Proposition 2.4.** Let  $\mathcal{M} = (\mathcal{Q}, \mathcal{I})$  be a  $\Pi$ -partite matroid and let  $D = \Pi(\mathcal{I})$  be its associated discrete polymatroid. If  $\mathcal{M}$  is  $\mathbb{K}$ -representable, then so is D. In addition, if D is  $\mathbb{K}$ -representable, then  $\mathcal{M}$  is representable over some finite extension of  $\mathbb{K}$ .



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Generalizacions of Matroids (Supermatroids and Convex Geometries)

Dunstan, Ingleton, and Welsh introduced the concept of a *supermatroid* in 1972 as a generalization of the concept of an ordinary matroid and integral polymatroid.

Distributive supermatroids or poset matroids are supermatroids defined on distributive lattices or sets of order ideals of partially ordered sets (posets).

Cg Matroids have been defined and worked by a research team of the Kyoto University – Research Institute of Mathematical Sciences and the Central Institute of Economics and Mathematics of the Russian Academy of sciences since 2006.

The concept of a distributive supermatroid (or a poset matroid) can be generalized by considering a convex geometry, instead of a poset, as the underlying combinatorial structure on which we define a matroidal structure, which we call a *cg-matroid*.

Characterizations of cg-matroids is done by means of the exchange property for bases and the augmentation property for independent sets.

Strict cg-matroids will turn out to be exactly cg-matroids that are also supermatroids. In other words, strict cg-matroids are exactly supermatroids defined on the lattices of closed sets of convex geometries.

Closure Space: Let E be a nonempty finite set and  $\mathcal{F}$  be a family of subsets of E. The pair  $(E, \mathcal{F})$  is called a *closure space* on E if it satisfies the following two conditions:

(F0) 
$$\emptyset$$
,  $E \in \mathcal{F}$ .

(F1) 
$$X, Y \in \mathcal{F} \Longrightarrow X \cap Y \in \mathcal{F}$$
.

**Posets:** A (non-strict) **partial order** is a binary relation " $\leq$ " over a set *P* which is antysymmetric, transitive, and reflexive, i.e., for all a, b, and c in P, we have that:  $a \leq a$  (reflexivity) for all a in P; if  $a \leq b$  and  $b \leq a$  then a = b (antisymmetry); if  $a \leq b$  and  $b \leq c$  then  $a \leq c$  (transitivity).

In other words, a partial order is an antisymmetric preorder.

A set with a partial order is called a **partially ordered set** (also called a **poset** 



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#### **SuperMatroids:**

Let us now deal with supermatroids, A pair  $(\mathcal{S}, \mathcal{F})$  is a *supermatroid on*  $\mathcal{S}$  if  $\mathcal{S} = (\mathcal{S}, \leqslant)$  is a poset with 0 and height function  $|\cdot|$ , and  $\mathcal{F} \subseteq \mathcal{S}$  is a subset such that:

(11)<sup>SM</sup> For all 
$$I \in \mathcal{F}$$
, if  $K \in \mathcal{S}$  and  $K \leqslant I$ , then  $K \in \mathcal{F}$ ;

$$(I2)^{\text{SM}}$$
  $|I| = |K|$  for all maximal  $I, K \in [0, A] \cap \mathscr{F}$  and all  $A \in \mathscr{S}$ .

#### **Convex Geometry:**

The set E is called the *ground set* of the closure space  $(E, \mathcal{F})$ , and each member of  $\mathcal{F}$  is called a *closed set*. Moreover, we call the closure space  $(E, \mathcal{F})$  a *convex geometry* if it satisfies the following condition:

(F2) 
$$\forall X \in \mathcal{F} \setminus \{E\}, \ \exists e \in E \setminus X \colon \ X \cup \{e\} \in \mathcal{F}.$$

Condition (F2) is equivalent to the following chain condition:

(F2)' Every maximal chain  $\emptyset = X_0 \subset X_1 \subset \cdots \subset X_n = E$  in  $\mathcal{F}$  has length n = |E|.

#### **Closure Operator:**

Next we define an operator  $\tau: 2^E \to 2^E$  associated with the closure space  $(E, \mathcal{F})$ . For any  $X \in 2^E$  define

$$\tau(X) = \bigcap \{Y \in \mathcal{F} \mid X \subseteq Y\}.$$

That is,  $\tau(X)$  is the unique minimal closed set containing X. The operator  $\tau$  satisfies the following properties (cl0) $\sim$ (cl3):

(c10) 
$$\tau(\emptyset) = \emptyset$$
.

(c11) 
$$X \subseteq \tau(X)$$
 for  $X \in 2^E$  (Extensionality).

(c12) 
$$X \subseteq Y \Longrightarrow \tau(X) \subseteq \tau(Y)$$
 for any  $X, Y \in 2^E$  (Monotonicity).

$$\text{(c13)} \quad \tau(\tau(X)) = \tau(X) \quad \text{ for any } X \in 2^E \quad \text{(Idempotence)}.$$

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In general, any operator  $\tau: 2^E \to 2^E$  satisfying the four conditions given above is called a *closure operator*. Conversely, given a closure operator  $\tau$ , define  $\mathcal{F} = \{X \in 2^E \mid \tau(X) = X\}$ . Then  $\mathcal{F}$  forms a closure space on E. Hence, for a finite set E and a closure operator  $\tau$  on E we also call the pair  $(E,\tau)$  a *closure space*.

In terms of closure operator, a closure space  $(E, \tau)$  is a convex geometry if and only if it satisfies the following property, called the *anti-exchange property*:

(AE) 
$$X \subseteq E, \ p \in E \setminus \tau(X), \ q \in \tau(X \cup \{p\}) \setminus \{p\} \implies p \notin \tau(X \cup \{q\}).$$

#### **Graded lattices Operations:**

Every convex geometry forms a graded lattice with respect to set-inclusion, where the lattice operations  $join \lor$  and  $meet \land$  are given by

$$X \vee Y = \tau(X \cup Y), \qquad X \wedge Y = X \cap Y$$

for any  $X, Y \in \mathcal{F}$ .

#### **Convex Geometry Matroids:**

For a convex geometry  $(E, \mathcal{F})$  and a

family  $\mathcal{B} \subseteq \mathcal{F}$ , suppose that  $\mathcal{B}$  satisfies the following three conditions:

- (B0)  $\mathcal{B} \neq \emptyset$ .
- (B1)  $B_1, B_2 \in \mathcal{B}$ ,  $B_1 \subseteq B_2 \implies B_1 = B_2$ .
- (BM) (Middle Base Property) For any  $B_1, B_2 \in \mathcal{B}$  and  $X, Y \in \mathcal{F}$  with  $X \subseteq B_1, B_2 \subseteq Y$ , and  $X \subseteq Y$ , there exists  $B \in \mathcal{B}$  such that  $X \subseteq B \subseteq Y$ .

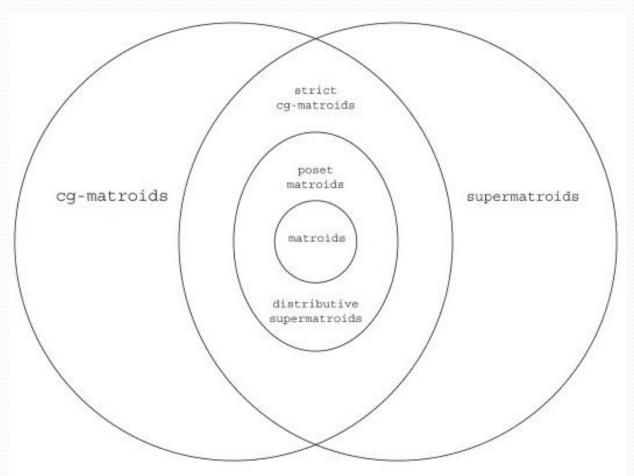
Then we call  $(E, \mathcal{F}; \mathcal{B})$  a matroid on the convex geometry  $(E, \mathcal{F})$  (or a cg-matroid for short). Each  $B \in \mathcal{B}$  is called a base, and  $\mathcal{B}$  the family of bases of cg-matroid  $(E, \mathcal{F}; \mathcal{B})$ .

Note that a cg-matroid  $(E, \mathcal{F}; \mathcal{B})$  is an ordinary matroid when  $\mathcal{F} = 2^E$  and that  $(E, \mathcal{F}; \mathcal{B})$  is a poset matroid (or a distributive supermatroid) when  $\mathcal{F}$  is the set of order ideals of a poset on E.



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#### Formal definition: Partition Representable Matroid

A matroid on the ground set N with the rank function r is said to be partition representable of degree  $d \ge 2$  if partitions  $\xi_i$ ,  $i \in N$ , of a finite set  $\Omega$  of the cardinality  $d^{r(N)}$ , exist such that the meet-partition  $\xi_I = \bigwedge_{i \in I} \xi_i$  has  $d^{r(I)}$  blocks of the same cardinality for every  $I \subset N$ .

#### Alternative Names: Secret Sharing Schemes and Almost Affine Codes

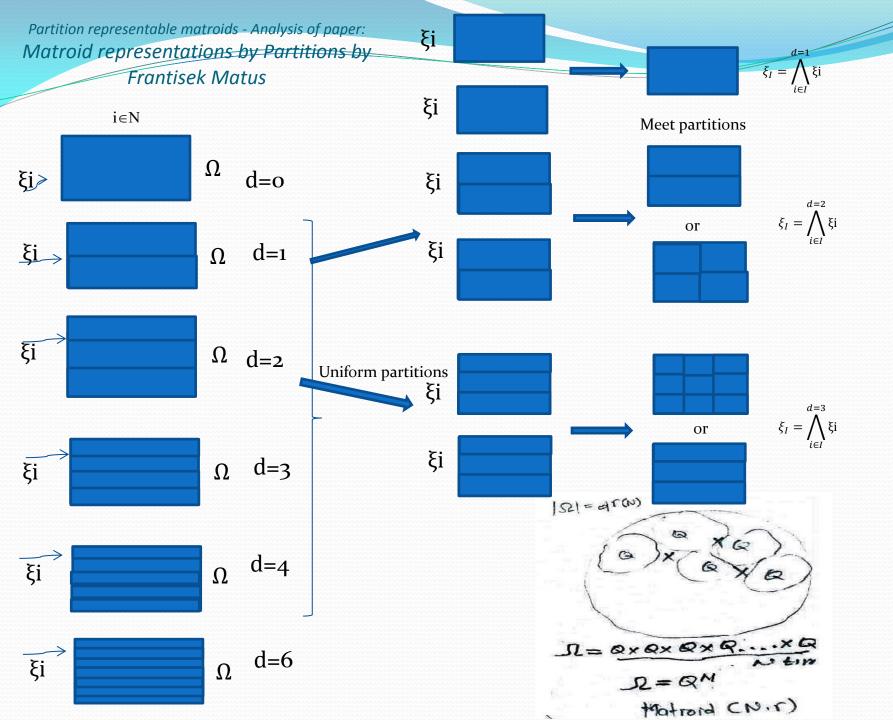
Partition representable matroids are called also secret-sharing or almost affinely representable and partition representations correspond to ideal secret-sharing schemes or to almost affine codes.

#### Meet Partitions of P-representation are Uniform:

A partition with all blocks having the same number of elements will be termed uniform. All meet-partitions of a p-representation are uniform:

All blocks of  $\xi_I$  have the cardinality  $d^{r(N)-r(I)}$ ; especially, ;  $\xi_{\varphi}$  has only one block being the whole set and  $\xi_N$  N has  $d^{r(N)}$  blocks being the singletons of  $\Omega$ .







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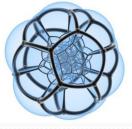
#### Motivation of the concept in Matroid Theory:

A central notion of the present study has emerged to us, maybe quite unexpectedly in context of the matroid representation theories, from an investigation of conditional independence structures among random variables, see [33]. After recognition of the role of entropy functions, cf. [17], some conditional independence structures were related to matroids and some of their probabilistic representations appeared to have a highly symmetric combinatorial form. The resulting notion is that of partition representable matroid

#### Other Research fields that have found the same concept:

was introduced independently in cryptology under the label 'secret-sharing matroid' [39] and in coding theory as 'almost affinely representable matroid' [40]. There, a starting point was the result of [8] relating ideal secret-sharing schemes to matroids.





Matroid representations by Partitions by Frantisek Matus

#### Formal definitions: Partitions are Random Variables

Suppose that  $\xi = (\xi_i)_{i \in N}$  is a system of partitions of a finite set  $\Omega$  and p is a nonnegative function on  $\Omega$  summing to one. Equivalently, we will say that  $\xi$  is a system of random variables.

 $\xi_I$ ,  $I \subset N$ , can be considered as a **Set Algebra** with **Probability space**  $(\Omega, p)$ 

Where the **Shannon entropy** is defined as:  $h_{\xi}(I) = -\sum_{A \in \xi_I} p(A) \ln p(A)$ 

And the Probability of the block A from  $\zeta_I$  is given by :  $p(A) = \sum_{\omega \in A} p(\omega)$ 

 $h_{\xi}$  is the **entropy function** of  $\xi$ .

 $(N, h_{\xi})$  is known to be a **Polymatroid**.



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What is the meaning that the  $\xi_I$ ,  $I \subset N$ , must be a set algebra? That they must obey the following properties:

#### The fundamental laws of set algebra

The binary operations of set union and intersection satisfy many identities.

PROPOSITION 1: For any sets A, B, and C, the following identities hold:

commutative laws:

- $A \cup B = B \cup A$
- $A \cap B = B \cap A$

associative laws:

- $(A \cup B) \cup C = A \cup (B \cup C)$
- $\bullet (A \cap B) \cap C = A \cap (B \cap C)$

distributive laws:

- $A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$   $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$

PROPOSITION 2: For any subset A of universal set U, where Ø is the empty set, the following identities hold:

identity laws:

- $\bullet A \cup \emptyset = A$
- $A \cap U = A$

complement laws:

- $A \cup A^C = U$
- $A \cap A^C = \emptyset$





#### Matroid representations by Partitions by Frantisek Matus

#### **Probability space of the Partitions:**

At the same time the  $\xi_I$ ,  $I \subset N$ , constitute a Probability space:

A probability space is a mathematical triplet  $(\Omega, \mathcal{F}, P)$  that presents a model for a particular class of real-world situations.

A probability space consists of three parts:

- A sample space, Ω, which is the set of all possible outcomes.
- A set of events F, where each event is a set containing zero or more outcomes.
- The assignment of probabilities to the events; that is, a function P from events to probabilities.

The prominent Soviet mathematician Andrey Kolmogorov introduced the notion of probability space, together with other axioms of probability, in the 1930s. An outcome is the result of a single execution of the model. Since individual outcomes might be of little practical use, more complex *events* are used to characterize groups of outcomes. The collection of all such events is a  $\sigma$ -algebra  $\mathcal{F}$ .

Finally, there is a need to specify each event's likelihood of happening. This is done using the probability measure function, P.

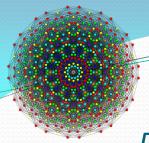
In short, a probability space is a measure space such that the measure of the whole space is equal to one.

The expanded definition is following: a probability space is a triple  $(\Omega, \mathcal{F}, P)$  consisting of:

- the sample space Ω an arbitrary non-empty set,
- the  $\sigma$ -algebra  $\mathcal{F} \subseteq 2^{\Omega}$  (also called  $\sigma$ -field) a set of subsets of  $\Omega$ , called events, such that:
  - F contains the empty set: ∅∈F,
  - F is closed under complements: if A∈F, then also (Ω\A)∈F,
  - F is closed under countable unions: if A<sub>i</sub>∈F for i=1,2,..., then also (∪<sub>i</sub>A<sub>i</sub>)∈F
    - The corollary from the previous two properties and De Morgan's law is that F is also closed under countable intersections: if A<sub>i</sub>∈F for i=1,2,..., then also (∩<sub>i</sub>A<sub>i</sub>)∈F
- the probability measure P: F→[0,1] a function on F such that:
  - P is countably additive: if {A<sub>i</sub>}⊆F is a countable collection of pairwise disjoint sets, then P(\(\preciou A\_i\)) = \(\sum P(A\_i)\), where "\(\preciou\)" denotes the disjoint union,
  - the measure of entire sample space is equal to one: P(Ω) = 1.



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#### Strongly probabilistically representable Matroids

Given this partition, if  $p(\omega) = |\Omega|^{-1}$  for  $\omega \in \Omega$  then  $h_{\xi} = r \cdot \ln d$  where d is the degree of  $\xi$ 

We say (N, r) to be *strongly probabilistically representable* if for a system  $\xi$  of random variables on  $(\Omega, p)$  the entropy function of  $\xi$  is proportional to the rank function r

$$h_{\xi} = c \cdot r$$

Every strong probabilistic representation  $\xi$  (over some  $(\Omega, p)$ ) of a connected matroid with rank at least two has associated a unique integer number  $d \ge 2$  such that  $p(A) = d^{-r(I)}$ ,  $A \in \xi_I$   $I \subset N$ 

A matroid is partition representable if and only if it is strongly probabilistically representable and that there are only trivial differences between partition representations and probabilistic representations.





Matroid representations by Partitions by Frantisek Matus

#### **Semimatroids** (Formal definition)

Let (N, g) be a polymatroid and |[g]| be the family of those triples (i, j | K),  $K \subset N$ , i,  $j \in N - K$ , which satisfy

$$g(i \cup K) + g(j \cup K) = g(K) + g(i \cup j \cup K).$$

The family [g] defined in this way is called Semimatroid.

#### Probabilistically representable Semimatroids:

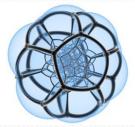
if for a system of random variables  $\xi$  on  $(\Omega, p)$  a triple (i, j | K) belongs to it if and only if  $\xi_i$  is stochastically conditionally independent of  $\xi_j$  given  $\xi_K$ 

A necessary and sufficient condition for this is  $h_{\xi}(i \cup K) + h_{\xi}(j \cup K) = h_{\xi}(K) + h_{\xi}(i \cup j \cup K)$  if i = j this means a functional dependence.

Thus, a semimatroid |[g]| is probabilistically representable if and only if  $|[g]| = |[h_{\xi}]|$  for some  $\xi$ 







#### Weakly probabilistically representable Matroids:

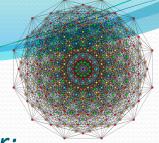
A matroid (N,r) is weakly probabilistically representable if the semimatroid [r] is probabilistically representable, i.e. if for a system of random variables on some  $(\Omega, p)$  the equality  $h_{\xi}(i \cup K) + h_{\xi}(j \cup K) = h_{\xi}(K) + h_{\xi}(i \cup j \cup K)$  is equivalent to  $r(i \cup K) + r(j \cup K) = r(K) + r(i \cup j \cup K)$  whatever  $K \subset N$  and  $i, j \in N - K$ .

#### **Secret-sharing schemes**

Let N be a finite set (of participants),  $0 \in N$  be a distinguished element of N (dealer), and  $\mathscr{A}$  a nonempty family (access structure) of subsets of N-0 no two of them being in a set inclusion. Let us suppose that every  $i \in N-0$  is contained in some  $J \in \mathscr{A}$ 

A system of random variables  $\xi = (\xi_i)_{i \in N}$  is called (probabilistic) secret-sharing scheme for  $\mathscr{A}$  if  $h_{\xi}(0 \cup I) = h_{\xi}(I)$  for every  $I \subset N - 0$  containing some  $J \in \mathscr{A}$ .





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#### Perfect secret-sharing scheme

A secret-sharing scheme is *perfect* if  $h_{\xi}(0 \cup I) = h_{\xi}(0) + h_{\xi}(I)$  for every  $I \subset N - 0$  containing no  $J \in \mathcal{A}$ ; in such a scheme if  $i \in N - 0$  then  $i \in J$  for some  $J \in \mathcal{A}$  and

$$h_{\xi}(i) \geqslant h_{\xi}(J) - h_{\xi}(J-i) = h_{\xi}(0 \cup J) - h_{\xi}(J-i)$$
  
 $\geqslant h_{\xi}((0 \cup J) - i) - h_{\xi}(J-i) = h_{\xi}(0).$ 

#### Ideal secret-sharing scheme

if 
$$h_{\xi}(i) = h_{\xi}(0) > 0$$
 for all  $i \in N - 0$ .

If  $\xi$  is an ideal secret-sharing scheme for some  $\mathscr A$  then there exists a unique connected matroid (N,r) such that

$$\mathscr{A} = \{J \subset N - 0; \ 0 \cup J \text{ is a circuit of } (N,r)\}$$
  $h_{\xi} = c \cdot r \text{ for a positive number } c$ 

Thus an ideal secret-sharing scheme is one of the strong probabilistic representations of a connected matroid and, omitting trivial situations, it is a partition representation of a connected matroid. In the reverse direction, all strong probabilistic representations give rise in a straightforward manner to ideal secret-sharing schemes.





Matroid representations by Partitions by Frantisek Matus

#### **Duality**

It is interesting to ask whether the dual of a p-representable matroid is p-representable?

There is a natural extension of the matroid duality to the class of semimatroids containing only the triples (i; j|K) with i /= j, see [33].

In [34] an example of a probabilistically representable semimatroid having the probabilistically nonrepresentable dual was found. Second, the p-representations of a matroid can have different structure than the p-representations of its dual.

For duality in perfect secret-sharing schemes see [21] and for duality in almost affine codes [40].



# Partition representable matroids - Analysis of paper: Matroid representations by Partitions by Frantisek Matus

# Comparing matroid representations

#### **Linearly representable Matroids:**

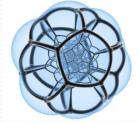
Linearly (l-) representable matroids are p-representable. If a matroid (N; r) is l-representable over a field then it is l-representable also over a finite field of sufficiently high cardinality d, that is, there are linear functionals  $x_i$ ,  $i \in N$ , on a linear space E of the dimension r(N) over the finite field such that r(I) is the rank of the hull of  $(x_i)$  if  $I \subset N$ .

Blocks of a partition  $\xi_i$  of  $\Omega = E$  are then taken as the shifts of the null-space of the functional  $x_i$ ,  $i \in N$  and it is almost straightforward that the system of partitions  $\xi = (\xi_i)_{i \in N}$  is a p-representation of the matroid (N; r) of degree d.

#### Multilinearly representable Matroid:

A matroid (N,r) is *multilinearly* (ml-) *representable* if at least one of the polymatroids (N,nr),  $n \ge 1$ , has an 1-representable expansion; equivalently, if for some  $n \ge 1$  there exist subspaces  $D_i$ ,  $i \in N$ , of a linear space E over a field such that the rank of the hull of  $D_i$ ,  $i \in I$ , is nr(I), for every  $I \subset N$ . If the subspaces have ranks at most 1 we have an 1-representation. Obviously, m1-representable matroids are p-representable.





Matroid representations by Partitions by Frantisek Matus

Obviously, ml-representable matroids are p-representable. The problem whether there exists a p-representable non-ml-representable matroid, remains open.

To solve, it would suffice to find a p-representable matroid such that its rank function does not satisfy the Ingleton inequality. This inequality is a necessary condition for a matroid to be ml-representable but entropy functions do exist violating it

An algebraically (a-) representable matroid need not be p-representable. A natural question is whether a p-representable and non-a-representable matroid exists.

Any minor of a matroid having a p-representation of some degree d is p-representable of the same degree d, [33,40].

A matroid is p-representable of degree two and three if and only if it is binary and ternary, respectively, see [33,5]. Even more, all its p-representations are p-isotopic to l-representations, [40].

Since the matroids  $L_n$ ,  $n \ge 2$  nonprime are non-p-representable and all their proper minors are l-representable, see [26,10], they are forbidden minors for the p-representability. Hence, the class of p-representable matroids has an infinite number of forbidden minors.



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#### The p-characteristic set of a matroid

It is the set of all integers  $d \ge 2$ ch that the matroid has a p-representation of degree d, cf. [20,23,18,38]. A natural question is which subsets of {2; 3; 4;:::} are the p-characteristic sets of matroids. We want to remark that the p-characteristic sets are multiplication closed.

In fact, if  $\xi$  and  $\eta$  are two p-representations of a matroid (N,r) having the degrees  $d_{\xi}$  and  $d_{\eta}$  and living on  $\Omega_{\xi}$  and  $\Omega_{\eta}$ , respectively, then their *product*  $\xi \otimes \eta = (\zeta_i)_{i \in N}$  will consist of the partitions  $\zeta_i$  of  $\Omega_{\xi} \times \Omega_{\eta}$  having the blocks  $A \times B$  for  $A \in \xi_i$  and  $B \in \eta_i$ . The product  $\xi \otimes \eta$  is obviously a p-representation of the matroid of degree  $d_{\xi} \cdot d_{\eta}$ .

Classifying the p-representations of a matroid, the most interesting are irreducible p-representations, i.e. those that are not p-isotopic to the product of two p-representations of the matroid. All the remaining p-representations are simply  $\xi \otimes \xi$ ,  $\xi \otimes \xi \otimes \xi$ , ....



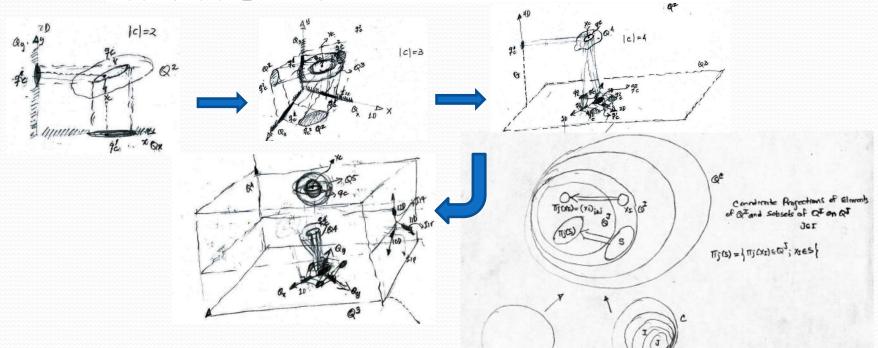


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#### Essential concepts involved in the Partition Representation of Matroids

#### **Coordinate Projection:**

Where  $Q \neq \emptyset$  and I are two sets, elements of  $Q^I$  will be denoted by  $x_I = (x_i)_{i \in I}$  and the coordinate projection of  $x_I \in Q^I$  on  $Q^J$ ,  $J \subset I$ , by  $\pi_J(x_I) = (x_i)_{i \in J}$ ;  $Q^\emptyset$  is a fixed one-element set. Subsets S of  $Q^I$  will be projected similarly  $\pi_J(S) = \{\pi_J(x_I) \in Q^J; x_I \in S\}$ .



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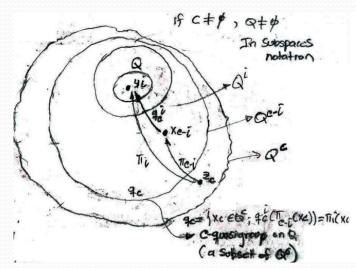
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#### Essential concepts involved in the Partition Representation of Matroids

#### Parastrophic quasigroups mappings:

 $\mathfrak{q}_C^i:Q^{C-i}\to Q$  are parastrophic quasigroups on Q (of arity |C|-1) in the usual sense; for background on the quasigroup theory see [14,15]. Note that for all  $i\in C$ 

$$\mathfrak{q}_C = \{x_C \in Q^C; \mathfrak{q}_C^i(\pi_{C-i}(x_C)) = \pi_i(x_C)\}.$$



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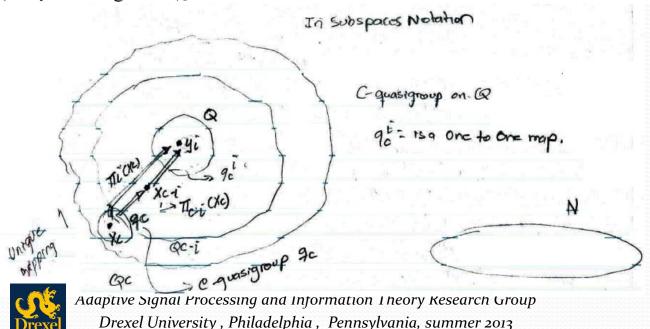


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#### Essential concepts involved in the Partition Representation of Matroids

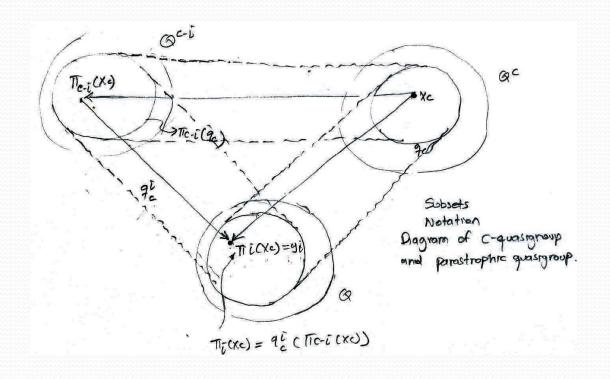
#### **C-Quasigroup:**

If C is a nonempty set then a C-quasigroup on  $Q \neq \emptyset$  will be for us a special subset  $\mathfrak{q}_C$  of  $Q^C$  defined as follows. For every  $i \in C$ , to every  $x_{C-i} \in Q^{C-i}$  there exists a unique  $y_i \in Q$  such that the C-tuple  $z_C$  composed from  $\pi_{C-i}(z_C) = x_{C-i}$  and  $\pi_i(z_C) = y_i$  belongs to  $\mathfrak{q}_C$ .



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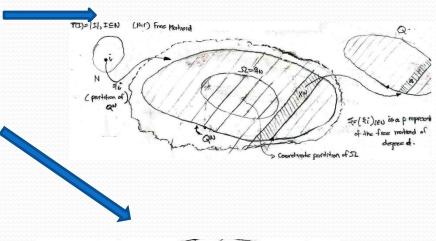
Review Chart: Parastrophic Quasigroups, C-Quasigroups, Coordinate Projections

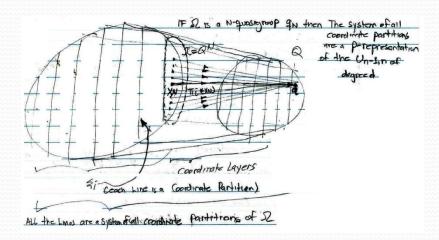


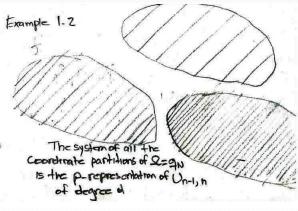


#### Matroid representations by Partitions by Frantisek Matus

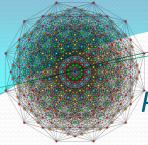
**Example 1.2.** Let (N,r) be the free matroid on N  $(r(I) = |I|, I \subset N)$  and Q a finite set of cardinality  $d \geqslant 2$ . We set  $\Omega = Q^N$  and  $\xi_i$ ,  $i \in N$ , to be the partition of  $\Omega$  with the blocks  $\{x_N \in \Omega; \pi_i(x_N) = a\}$ ,  $a \in Q$ . Verbally expressed, the blocks are ith parallel coordinate layers of  $Q^N$ . It is obvious that  $\xi = (\xi_i)_{i \in N}$  is a p-representation of the free matroid of degree d. If a partition of an arbitrary subset  $\Omega$  of  $Q^N$  is specified by the coordinate projection as above (its blocks are intersections of the coordinate layers with  $\Omega$ ) we speak about a *coordinate partition* of  $\Omega$ . The system of all coordinate partitions of any N-quasigroup  $\Omega = \mathsf{q}_N \subset Q^N$  on Q can be easily identified as a p-representation of the uniform matroid  $\mathsf{U}_{n-1,n}$  of degree d. This matroid has a ground set N of the cardinality  $n \geqslant 1$  and the rank function  $r(I) = \min\{|I|, n-1\}$ ,  $I \subset N$ ; the only circuit of the matroid is C = N.











d rest

# Partition representable matroids - Analysis of paper:

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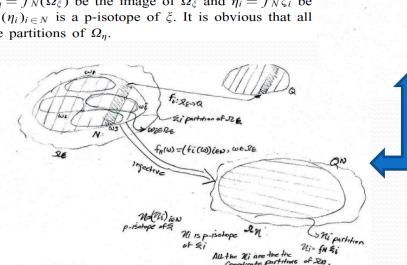
#### Essential concepts involved in the Partition Representation of Matroids

#### P-isotopic representations:

**Definition** Two p-representations  $\xi = (\xi_i)_{i \in N}$  and  $\eta = (\eta_i)_{i \in N}$  of a matroid (N, r), living on two sets  $\Omega_{\xi}$  and  $\Omega_{\eta}$ , respectively, are *p-isotopic* if there exists a bijection f of  $\Omega_{\xi}$  onto  $\Omega_{\eta}$  such that  $f \xi_i = \eta_i$  for all  $i \in N$ .

P-isotopic p-representations of a matroid must be of the same degree;  $\eta$  is said to be a *p-isotope* of  $\xi$ . Blocks of  $\xi_I$  are mapped by f onto blocks of  $\eta_I$  for all  $I \subset N$ .

Let  $\xi$  be a p-representation of a matroid of degree d living on  $\Omega_{\xi}$ . If Q is a set of the cardinality d and  $f_i: \Omega_{\xi} \to Q$  is a function distinguishing and constant on the blocks of  $\xi_i$ ,  $i \in N$ , then the composed function  $f_N$  defined by  $f_N(\omega) = (f_i(\omega))_{i \in N}$ ,  $\omega \in \Omega_{\xi}$ , maps  $\Omega_{\xi}$  injectively into  $Q^N$ . Let  $\Omega_{\eta} = f_N(\Omega_{\xi})$  be the image of  $\Omega_{\xi}$  and  $\eta_i = f_N \xi_i$  be the partitions of  $\Omega_{\eta}$ . Obviously,  $\eta = (\eta_i)_{i \in N}$  is a p-isotope of  $\xi$ . It is obvious that all partitions within  $\eta$  are the coordinate partitions of  $\Omega_{\eta}$ .

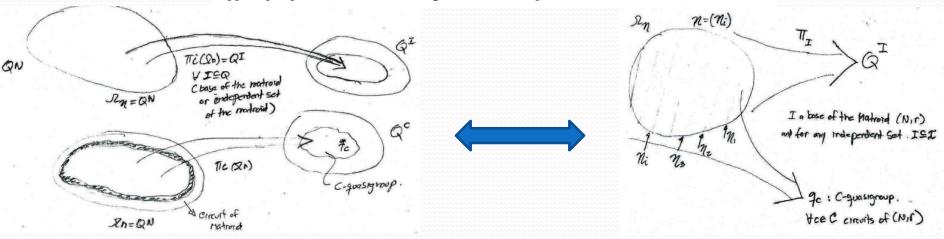


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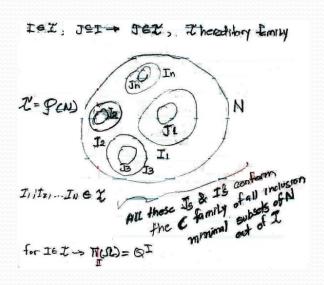
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#### Essential concepts involved in the Partition Representation of Matroids

We have just seen that every p-representation  $\xi$  of a matroid is p-isotopic to the p-representation  $\eta$  which lives on a subset of a Cartesian product and consists of all coordinate partitions of the subset. It is not difficult to see that  $\pi_I(\Omega_\eta) = Q^I$  for any base I of the matroid (and consequently for every independent set of the matroid) and that  $\pi_C(\Omega_\eta)$  is a C-quasigroup for every circuit C of the matroid. The following lemma asserts that these two properties characterize those subsets of Cartesian products that support p-representations consisting of coordinate partitions.

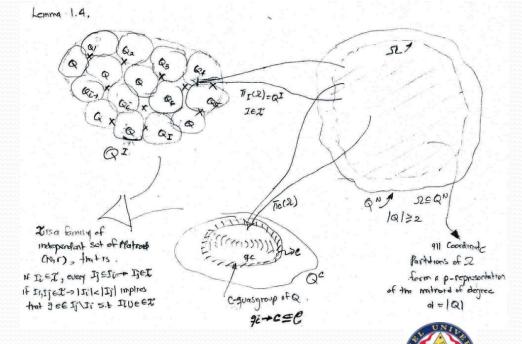


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**Lemma 1.4.** Let N be a finite nonempty set and  $\mathcal{I}$  a nonempty hereditary family of subsets of N ( $I \in \mathcal{I}$  and  $J \subset I$  implies  $J \in \mathcal{I}$ ). We denote by  $\mathscr{C}$  the family of all inclusion-minimal subsets of N out of  $\mathcal{I}$ . If  $\Omega$  is a subset of a Cartesian product  $Q^N$ ,  $|Q| \ge 2$  finite, such that  $\pi_I(\Omega) = Q^I$  for  $I \in \mathcal{I}$  and  $\pi_C(\Omega)$  is a C-quasigroup

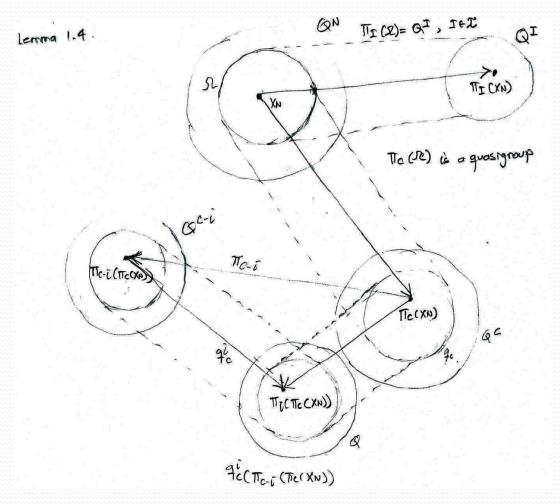
on Q for  $C \in \mathcal{C}$  then  $\mathcal{I}$  is the family of independent sets of a matroid on N. The coordinate partitions of the set  $\Omega$  form a p-representation of the matroid of the degree |Q|.



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Review Chart: Concepts already defined and involved in lemma 1.4





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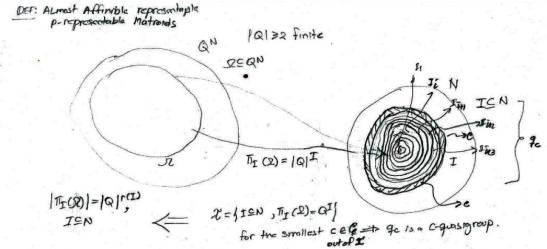


Matroid representations by Partitions by Frantisek Matus

#### Essential concepts involved in the Partition Representation of Matroids

Almost Affine codes & Almost affinely representable matroids.

In [40], a subset  $\Omega$  of  $Q^N$ ,  $|Q| \ge 2$  finite, is called *almost affine code* if the cardinality of every projection  $\pi_I(\Omega)$ ,  $I \subset N$ , is a power of |Q|. The family  $\mathscr{I} = \{I \subset N; \pi_I(\Omega) = Q^I\}$  is nonempty, hereditary and obviously for the minimal C out of  $\mathscr{I}$  the projection  $\pi_C(\Omega)$  is a C-quasigroup. The above lemma implies that  $|\pi_I(\Omega)| = |Q|^{r(I)}$ ,  $I \subset N$ , for a unique matroid (N, r). The matroids arising in this way, for us the p-representable matroids, are called in [40] almost affinely representable.





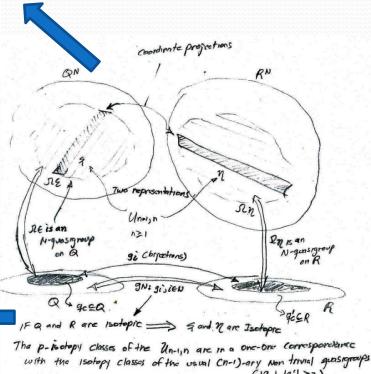
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#### Expansion of the the example 1.2

It is trivial that for the free matroid any two of its p-representations of the same degree are p-isotopic. Let us suppose now that  $\xi$  and  $\eta$  are two p-representations of  $\bigcup_{n=1,n}, n \ge 1$ , living on  $\Omega_{\xi} \subset Q^N$  and  $\Omega_{\eta} \subset R^N$ and consisting of the coordinate projections, correspondingly. Necessarily,  $\Omega_{\varepsilon}$  is an N-quasigroup on Q and  $\Omega_n$  is an N-quasigroup on R. If the two N-quasigroups are isotopic, i.e. there exist bijections  $g_i: Q \to R$ ,  $i \in N$ , such that the composed mapping

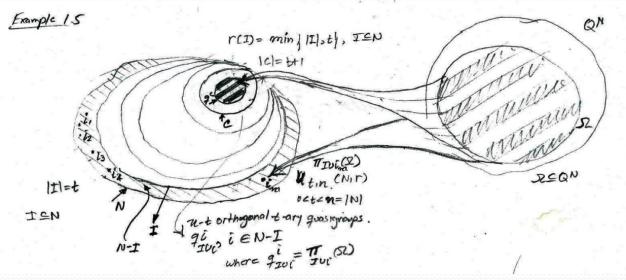
 $g_N$  is a bijection between  $\Omega_{\xi}$  and  $\Omega_{\eta}$ , then  $\xi$  and  $\eta$  are obviously p-isotopic. But it is not difficult to reverse the assertion, namely, if  $\xi$  and  $\eta$  are p-isotopic then every p-isotopy must have the composed form  $g_N$  for some bijections  $g_i$ ,  $i \in N$ . Hence, the p-isotopy classes of the p-representations of  $U_{n-1,n}$  are in a one-to-one correspondence with the isotopy classes of the usual (n-1)-ary nontrivial finite quasigroups.





#### Matroid representations by Partitions by Frantisek Matus

**Example 1.5.** Let (N,r) be the uniform matroid  $U_{t,n}$  where 0 < t < n = |N| and  $r(I) = \min\{|I|,t\}$ ,  $I \subset N$ . A p-representation of  $U_{t,n}$  is nothing else but a set  $\Omega \subset Q^N$  such that the projections  $\pi_C(\Omega)$  are C-quasigroups for all C of the cardinality t+1; the condition on bases from Lemma 1.4. follows and is thus redundant. These sets  $\Omega$  correspond, by a fixed base I, |I| = t, to the families of n-t orthogonal t-ary quasigroups  $\mathfrak{q}_{I \cup i}^i$ ,  $i \in N-I$ , where  $\mathfrak{q}_{I \cup i} = \pi_{I \cup i}(\Omega)$ . (The orthogonality can be defined by demanding injectiveness of the mapping  $x_I \mapsto (\pi_{I \cap J}(x_I), (\mathfrak{q}_{I \cup j}^j(x_I))_{j \in J-I})$  for every  $J \subset N$ , |J| = t.) In a combinatorial language, the p-representations of uniform matroids of rank at least two are the orthogonal latin hypercubes, see [14]. In coding theory one uses, equivalently, the term 'orthogonal arrays of size  $|Q|^t$ , n constraints, |Q| levels, strength t, and index 1',





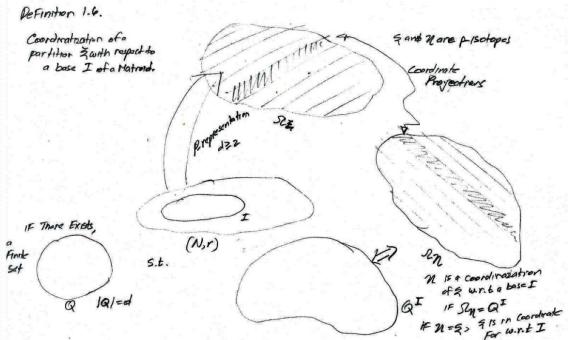
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Matroid representations by Partitions by Frantisek Matus

#### Essential concepts involved in the Partition Representation of Matroids

#### P-isotope Coordinatization

**Definition 1.6.** Let (N,r) be a matroid and  $\xi$  its p-representation of degree  $d \ge 2$  on  $\Omega_{\xi}$ . A p-isotope  $\eta$  of  $\xi$  living on  $\Omega_{\eta}$  is called *coordinatization of*  $\xi$  *w.r.t.* a base I of the matroid if  $\Omega_{\eta} = Q^I$  for some finite set Q of cardinality d and  $\eta_i$ ,  $i \in I$ , are the coordinate partitions of  $Q^I$ . When  $\eta = \xi$  here we say that  $\xi$  is in the *coordinate form* w.r.t. I.





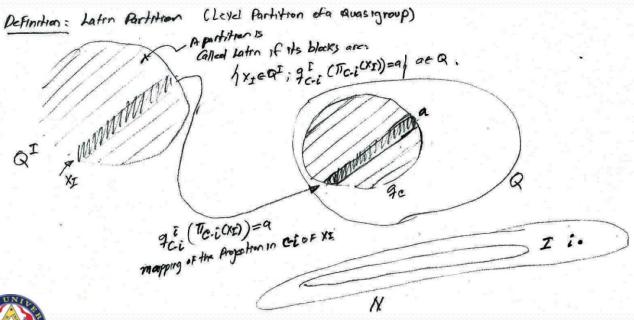
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#### Essential concepts involved in the Partition Representation of Matroids

#### Latin Partitions (Level partitions of Quasigroups)

It is easy to see that every p-representation  $\xi$  can be brought into the coordinate form w.r.t. any basis I of the underlying matroid; it suffices to transform it by a composed function  $f_I$ . A partition of  $Q^I$  is called a *latin partition* if its blocks can be given as  $\{x_I \in Q^I; \mathfrak{q}_{C-i}^i(\pi_{C-i}(x_I)) = a\}$ ,  $a \in Q$ , where  $\mathfrak{q}_C$  is a C-quasigroup on Q,  $i \notin I$ , and  $i \in C \subset i \cup I$ . We will use equivalently also the term *level partition of a quasigroup*. Latin partitions are obviously uniform.



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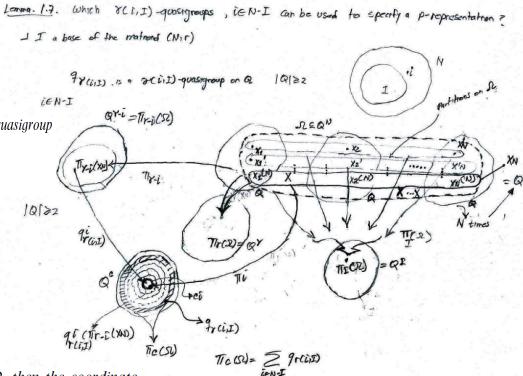
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#### Essential concepts involved in the Partition Representation of Matroids

All 'nonbasical' partitions of a p-representation in the coordinate form w.r.t. a base I are latin, they are constructed from  $\gamma(i,I)$ -quasigroups,  $i \in N - I$ . To recognize which  $\gamma(i,I)$ -quasigroups,  $i \in N - I$ , can be used to specify a p-representation we will use later this lemma.

**Lemma 1.7.** Let I be a base of a matroid (N,r),  $\mathfrak{q}_{\gamma(i,I)}$ ,  $i \in N-I$ , be a  $\gamma(i,I)$ -quasigroup on Q,  $|Q| \ge 2$ , and

$$\Omega = \{x_N \in Q^N; \pi_i(x_N) = \mathfrak{q}^i_{\gamma(i,I)}(\pi_{\gamma(i,I)-i}(x_N)), \ i \in N-I\}.$$



If  $\pi_C(\Omega)$  is a C-quasigroup on Q for every circuit C,  $|C-I| \ge 2$ , then the coordinate partitions of  $\Omega$  (of  $Q^I$  together with the level partitions of  $\mathfrak{q}_{\gamma(i,I)}$ ,  $i \in N-I$ ) provide a p-representation of (N,r) (in the coordinate form w.r.t. I).



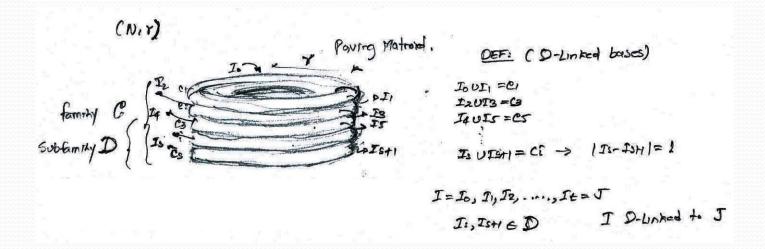


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#### Essential concepts involved in the Partition Representation of Matroids

#### D-Linked bases of a Matroid

Let  $\mathscr{D}$  be a subfamily of the family  $\mathscr{C}$  of all circuits of a matroid. Two bases I and J of the matroid are  $\mathscr{D}$ -linked if there exists a sequence of bases  $I = I_0, I_1, \dots, I_t = J$ ,  $t \ge 0$ , such that  $|I_s - I_{s+1}| = 1$  and the unique circuit contained in the union of  $I_s$  and  $I_{s+1}$  belongs to  $\mathscr{D}$ ,  $0 \le s < t$ . To be linked is a relation of equivalence.

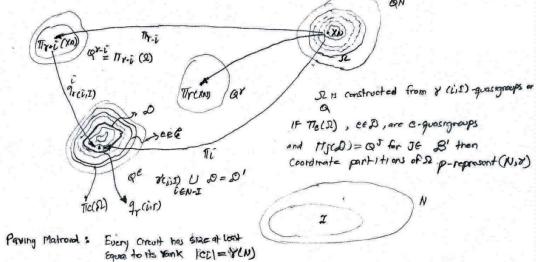




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**Lemma 1.8.** Let (N,r) be a paving matroid,  $r(N) \ge 2$ , I be its base and the family  $\mathscr{D}$  contain its circuits  $C, |C| = r(N), |C - I| \ge 2$ . Let  $\mathscr{D}' = \mathscr{D} \cup \{\gamma(i,I); i \in N - I\}$  and  $\mathscr{B}'$  be an arbitrary family of its bases such that every base not  $\mathscr{D}'$ -linked to I is  $\mathscr{D}'$ -linked to some base from  $\mathscr{B}'$ . Let a set  $\Omega$  be constructed from  $\gamma(i,I)$ -quasigroups on Q as in Lemma 1.7. If the projections  $\pi_C(\Omega)$  are C-quasigroups for  $C \in \mathscr{D}$  and  $\pi_J(\Omega) = Q^J$  for  $J \in \mathscr{B}'$  then the coordinate partitions of  $\Omega$  p-represent the matroid (N,r).

Lemma 1.9 (U,r) st.  $r(N) \ge 2$ , I'm its base,  $|C| = r(N) > |C-I| \ge 2$ I  $D' = \mathcal{D} \cup \{ \mathcal{T}(U,I) ; i \in N-I \}$ B is an arbitrary family of bases, Set if Si is not D-Linked to I is  $\mathcal{D}'$  Linked to  $\mathcal{D} \in \mathcal{B}'$ 





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#### Essential concepts involved in the Partition Representation of Matroids

### **Matroid Quasigroup Equations**

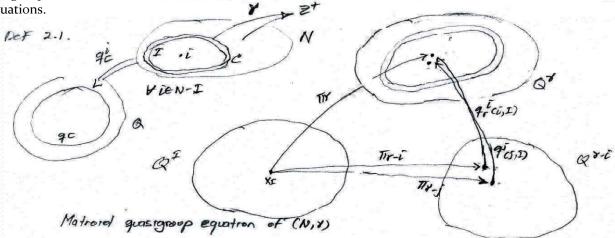


**Definition 2.1.** Let (N,r) be a matroid and Q a nonempty set. A family of C-quasi-groups  $\mathfrak{q}_C$  on Q for C running through all circuits of the matroid solves the *matroid* quasigroup equations if the following is satisfied. For every base I, for every  $i \in N-I$ , for every circuit C containing i and for every  $x_I \in Q^I$ 

$$\mathfrak{q}_{\gamma(i,I)}^i(\pi_{\gamma(i,I)-i}(x_I))=\mathfrak{q}_C^i(y_{C-i}),$$

where  $y_{C-i} = (y_j)_{j \in C-i}$  has the coordinates  $y_j = \mathfrak{q}^J_{\gamma(j,I)}(\pi_{\gamma(j,I)-j}(x_I))$  for  $j \in C - (i \cup I)$  and  $y_j = x_j$  for  $j \in C \cap I$ . In a single instance of I, i and C we speak about a *matroid* quasigroup equation of (N, r).

The unknowns in these equations are quasigroups (functions) which after substitutions of their arguments satisfy ordinary quasigroup equations.





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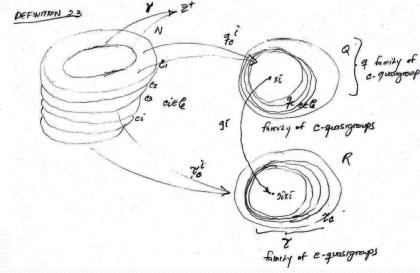
#### Essential concepts involved in the Partition Representation of Matroids

#### Simulateneously Isotopic family of C-quasigroups of a Matroid:

**Definition 2.3.** Let (N,r) be a matroid with the collection of circuits  $\mathscr{C}$ . A family  $(\mathfrak{q}_C)_{C\in\mathscr{C}}$  of C-quasigroups on a set Q is said to be *simultaneously isotopic* with a family  $(\mathfrak{r}_C)_{C\in\mathscr{C}}$  of C-quasigroups on a set R if there exist bijections  $g_i:Q\to R,\ i\in N,$  such that  $(x_i)_{i\in C}\in\mathfrak{q}_C$  if and only if  $(g_ix_i)_{i\in C}\in\mathfrak{r}_C$  for all  $C\in\mathscr{C}$ .

It is a matter of easy technicalities to see that if a family  $(q_C)_{C\in\mathscr{C}}$  solves the matroid quasigroup equations then its every simultaneous isotope solves the same set of equations as well. The simultaneous isotopy is a natural equivalence on the class of all solutions of the matroid quasigroup equations. The main result of this section establishes a connection between p-representations and matroid quasigroup equations.

**Proposition 2.4.** For any matroid (N,r) there is a one-to-one correspondence between the p-isotopy classes of the p-representations of (N,r) and the simultaneous isotopy classes of the nontrivial solutions of the matroid quasigroup equations associated to (N,r).







Essential concepts involved in the Partition Representation of Matroids

Necessary and sufficient condition for a Matroid to be P-representable

Now, it becomes evident that the p-representations can be systematically studied in the disguise of quasigroup equations: according to Proposition 2.4. a matroid is p-representable if and only if the system of its matroid quasigroup equations has a nontrivial solution. A look into the first part of the previous proof reveals that the equality  $\Omega = \Omega_{[I]}$  for a base I is equivalent to the assertion ' $(\mathfrak{q}_C)_{C \in \mathscr{C}}$  solves the subsystem of those matroid quasigroup equations which are indexed by the very base I, any  $i \in N - I$  and any circuit C containing i'. If the equality takes place for one base then it is valid for all bases because the sets  $\Omega_{[I]}$  contain  $\Omega$  and have the same cardinality. Therefore, while solving the matroid quasigroup equations for a matroid one can fix an appropriate base I of the matroid freely,



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#### Essential concepts involved in the Partition Representation of Matroids

Generalized Associativity of Quasigroups

**Consequence 3.2.** If four binary quasigroup operations on a finite set Q satisfy the equation  $(x \circ y) * z = x \star (y \diamond z)$  for all  $x, y, z \in Q$  then there exist a group  $(Q, \cdot)$ , unique up to isomorphisms, and five permutations  $\alpha, \beta, \gamma, \delta, \varepsilon$  of Q such that

$$x \circ y = \delta^{-1}(\alpha(x) \cdot \beta(y)), \quad x * z = \delta(x) \cdot \gamma(z),$$
  
 $y \diamondsuit z = \varepsilon^{-1}(\beta(y) \cdot \gamma(z)), \quad x \star z = \alpha(x) \cdot \varepsilon(z)$ 

is valid for all  $x, y, z \in Q$ .

**Consequence 3.4.** If six binary quasigroup operations on a finite set Q satisfy the equation  $(x \circ y) * (v \square u) = (x \bullet v) \star (y \blacksquare u)$  for all  $x, y, u, v \in Q$  then there exist an Abelian group, unique up to isomorphisms, and eight permutations  $\alpha, \beta, \gamma, \delta, \hat{\alpha}, \hat{\beta}, \hat{\gamma}$  and  $\hat{\delta}$  of Q such that

$$x \circ y = \hat{\alpha}^{-1}(\alpha(x) + \beta(y)), \quad x \bullet v = \hat{\delta}^{-1}(\alpha(x) + \delta(v)), \quad x * y = \hat{\alpha}(x) + \hat{\gamma}(y),$$
$$v \Box u = \hat{\gamma}^{-1}(\delta(v) + \gamma(u)), \quad y \bullet u = \hat{\beta}^{-1}(\beta(y) + \gamma(u)), \quad x \star y = \hat{\delta}(x) + \hat{\beta}(y).$$



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### Essential concepts involved in the Partition Representation of Matroids

#### **General Distributivity of Quasigroups**

The general distributive law  $x \star (y \circ z) = (x \Box y) * (x \diamondsuit z)$ ,  $x, y, z \in Q$ , has not been solved completely for arbitrary quasigroups, cf. [1,2,24]. It is not a balanced equation, see [41]. Obviously, when five quasigroup operations  $\star$ ,  $\circ$ ,  $\Box$ , \* and  $\diamondsuit$  solve the equation then also the following five quasigroup operations:

$$y \circledcirc z = \hat{\alpha}^{-1}(\beta(y) \circ \gamma(z)),$$

$$x \circledcirc y = \phi(\alpha(x) \bigstar \hat{\alpha}(y)), \quad x \circledcirc y = \hat{\beta}^{-1}(\alpha(x) \Box \beta(y)),$$

$$x \circledcirc y = \phi(\hat{\beta}(x) * \hat{\gamma}(y)), \quad x \circledcirc z = \hat{\gamma}^{-1}(\alpha(x) \diamondsuit \gamma(z)),$$

constructed by use of arbitrary seven permutations  $\alpha, \beta, \gamma, \hat{\alpha}, \hat{\beta}, \hat{\gamma}$  and  $\phi$  solve the equation, as well. The two quintuples of operations will be termed *equivalent*.

**Proposition 3.5.** Let Q be a finite set of cardinality  $d \ge 2$ . There is a bijective correspondence between the equivalence classes of the solutions  $(\star, \circ, \Box, *, \diamond)$  of the general distributive law, consisting of quintuples of quasigroups on Q, and the p-isotopy classes of degree d of the matroid  $P_7$  from Fig. 5 left.

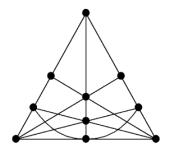


Matroid representations by Partitions by Frantisek Matus

#### Essential concepts involved in the Partition Representation of Matroids Non-p-representable matroids

**Proposition 4.1.** The two matroids of rank three from Fig. 7 are not p-representable.





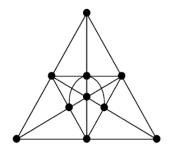
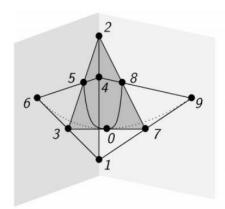


Fig. 7. Two examples of non-p-representable matroids.

**Proposition 4.2.** The two nonDesargues matroids (on ten element sets) of the ranks three and four are both non-p-representable.



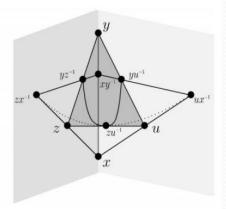




Fig. 9. NonDesargues matroids are non-p-representable.

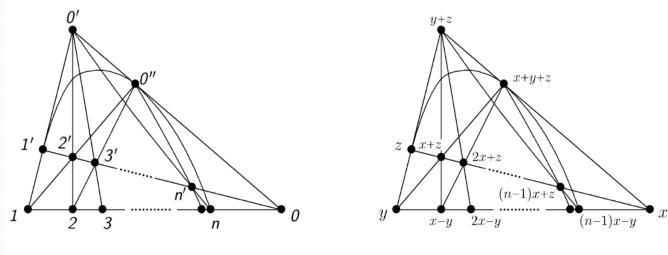
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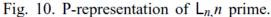
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Matroid representations by Partitions by Frantisek Matus

# Essential concepts involved in the Partition Representation of Matroids Non-p-representable matroids

**Proposition 4.3.** The matroid  $L_n$  is p-representable if and only if n is prime. In this case every p-representation of  $L_n$  is p-isotopic to the one given in Fig. 10 right where the addition is in the group  $\mathbb{Z}_n^m, m \ge 1$ .







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